

Lecture 10

Beyond Global Sensitivity Local Sensitivity

$$\Delta_{GS} = \max_{X, X'} \|f(X) - f(X')\|$$

$f: X^n \rightarrow \mathbb{R}$

$$\Delta_{LS}^{(f)}(X) = \max_{X' \text{ nbrs to } X} |f(X) - f(X')|$$

Before: $f(X) + \text{Lap}\left(\frac{\Delta_{GS}}{\varepsilon}\right)$ ✓ (2,0)-DP

Now: $f(X) + \text{Lap}\left(\frac{\Delta_{LS}(X)}{\varepsilon}\right) X$

$f(X)$: dist btw 2 closest points in $X \in \mathbb{R}^n$

$$\Delta_{GS} = 0 \quad X = \{0, 0\}, X' = \{0, \infty\}$$

$$X = \{0, 0, 0\}$$

$$X' = \{0, 0, \infty\}$$

$$\Delta_{LS}(X) = 0$$

$f(X) + \text{Lap}(0) = f(X) \Rightarrow 0$ deterministically

$$X = \{0, 0, 10000\}$$

$$\Delta_{LS}(X) = 10000$$

$$f(X') = \text{Lap}\left(10000/\varepsilon\right) = \text{Lap}\left(10000/\varepsilon\right)$$

Propose-Test-Release Dwork - Lei '09

1. Propose: bound on LS
2. Test: Is it?
3. Release: If yes, $f(x) + \text{Lap}\left(\frac{\text{LS}}{\epsilon}\right)$

- ϵ -DP
1. Propose a bound β on LS at x
 2. Compute $\gamma = \min_x d(x, \hat{x})$, where $\Delta_{LS}(\hat{x}) \geq \beta$ $\leftarrow |x|^{1/\epsilon}_{\text{slow}}$
 3. $\hat{x} = x + \text{Lap}(1/\epsilon)$ Hamming dist, # of pts to change
 4. If $\gamma \leq \ln(1/\delta)/\epsilon$, return \perp
 5. If $\gamma > \ln(1/\delta)/\epsilon$, return $f(x) + \text{Lap}(\beta/\epsilon)$.
- Compute LS: $|x|^\epsilon$

Theorem: PTR is $(2\epsilon, \delta)$ -DP

Proof: \perp

$$\Pr[M(x) = \perp] \in [e^{-\epsilon}, e^{\epsilon}] \cdot \Pr[M(x') = \perp]$$

Case 1: $\Delta_{LS}(x) > \beta$

$$\gamma = 0. \quad \Pr[\hat{x} > \ln(1/\delta)/\epsilon] \leq \delta$$

$\nwarrow \Pr[M(x) \neq \perp] \leq \delta$

$$T \subseteq \mathbb{R} \cup \{\perp\}$$

$$\begin{aligned} \Pr[M(x) \in T] &= \Pr[M(x) \in T \cap \{\perp\}] + \Pr[M(x) \in T \cap \mathbb{R}] \\ &\leq e^{\epsilon} \Pr[M(x') \in T \cap \{\perp\}] + \Pr[M(x) \neq \perp] \\ &\leq e^{\epsilon} \Pr[M(x') \in T] + \delta \end{aligned}$$

(ϵ, δ) -DP

Case 2: $\Delta_{LS}(X) \leq \beta$

1. Release \hat{Y} : ϵ -DP
2. Release $f(X) + \text{Lap}(\beta/\epsilon)$: ϵ -DP
 $(2\epsilon, 0)$ -DP \rightarrow

Application: Histograms (modal element)

Laplace Histogram: L_∞ err $\approx \log |X|$ error ϵ -DP

Relax to approx \rightarrow no dependence on $|X|$.

Problem: $X \in X^n$. Compute most freq elt.

$X = \{v, v, \dots, v\}$, how many pts x to change before mode changes?
 $\approx \frac{1}{2}(\text{count of most freq elt} - \text{count of 2nd most})$

1. Propose $LS \leq 0$.

2. $\hat{Y} =$

3. $\hat{Y} = Y + \text{Lap}(1/\epsilon)$. 4. If $\hat{Y} \leq \log(1/\delta)/\epsilon$, 1. 5. If $\hat{Y} \geq \log(1/\delta)/\epsilon$, output $\text{mode}(x)$.

$\Pr[\text{Lap}(1/\epsilon) > \ln(1/\delta)/\epsilon] \leq \delta$. Want: $\hat{Y} \geq 2\log(1/\delta)/\epsilon$.

Need: diff #1 and #2 elt $\geq 4\ln(1/\delta)/\epsilon$,

Theorem 3. There exists an (ϵ, δ) -differentially private algorithm which identifies the most frequent element from an arbitrary dataset with probability at least $1 - \delta$, as long as the gap between the count of the most frequent and the second most frequent element is at least $4\ln(1/\delta)/\epsilon$.

Thm 3.5 Vadhan '17

Theorem 4. There exists an (ϵ, δ) -differentially private algorithm which can, with high probability, output the count of every item in a dataset up to additive $O(\log(1/\delta)/\epsilon)$.

ϵ -DP $O(\log |X|/\epsilon)$

Privately Bounding Local Sensitivity

Before: Guessed LS bound

Now: Est LS. $\ln(1/\epsilon)$

1. Compute $\hat{\gamma} = \Delta_{LS}^{(A)}(x) + \text{Lap}\left(\frac{\Delta_{LS}^{(A)}}{\epsilon}\right) + \ln(1/\epsilon)/2$

2. Output $f(x) + \text{Lap}(\hat{\gamma}/\epsilon)$.

Smooth Sensitivity Nissim Raskhadnikova Smith '07

$$\Delta_{ss}^{(\epsilon)}(x) = \max_{\tilde{x} \in X^n} \left\{ \Delta_{ls}^{(\epsilon)}(\tilde{x}) e^{-\epsilon d(x, \tilde{x})} \right\}$$

↑ Hamming.

$$\cancel{f(x) + \text{Lap}\left(\frac{\Delta_{ss}(x)}{\epsilon}\right)} \quad f(x) + \text{Cauchy}\left(\frac{\Delta_{ss}(x)}{\epsilon}\right)$$

↑ (ϵ, δ) -DP



ϵ -DP