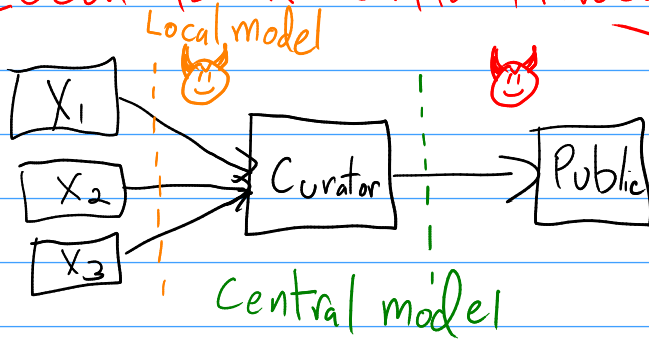


# Lecture 17

## Deployments of DP: Local Differential Privacy

### Local Differential Privacy



$$X_1, \dots, X_n \in \{0, 1\}, \mu = \frac{1}{n} \sum X_i$$

$$\hat{\mu} = \frac{1}{n} \sum X_i + \text{Lap}\left(\frac{1}{\epsilon n}\right)$$

$$|\hat{\mu} - \mu| \leq \frac{1}{\epsilon n} \iff \text{If want } |\hat{\mu} - \mu| \leq \alpha, n \geq \frac{1}{\alpha \epsilon}$$

Randomized Response  $\frac{1}{2} + \gamma, \frac{1}{2} - \gamma$

$$Y_i = \begin{cases} X_i & \text{w.p. } \frac{e^\epsilon}{1+e^\epsilon} \\ 1-X_i & \text{w.p. } \frac{1}{1+e^\epsilon} \end{cases} \quad \left(\frac{e^\epsilon}{1+e^\epsilon}\right) / \left(\frac{1}{1+e^\epsilon}\right) = e^\epsilon$$

$$\hat{\mu} = \frac{1}{n} \sum \left( \frac{e^\epsilon + 1}{e^\epsilon - 1} \cdot Y_i - \frac{1}{e^\epsilon - 1} \right)$$

$$E[\hat{\mu}] = \frac{1}{n} \sum E \left[ \frac{e^\epsilon + 1}{e^\epsilon - 1} \cdot Y_i - \frac{1}{e^\epsilon - 1} \right]$$

$$= \frac{1}{n} \sum \left( \frac{e^\epsilon + 1}{e^\epsilon - 1} \cdot \left( X_i \cdot \frac{e^\epsilon}{1+e^\epsilon} + (1-X_i) \cdot \frac{1}{1+e^\epsilon} \right) - \frac{1}{e^\epsilon - 1} \right)$$

$$= \frac{1}{n} \sum \left( \frac{1}{e^\epsilon - 1} - \frac{1}{e^\epsilon - 1} \frac{X_i \cdot e^\epsilon - X_i}{e^\epsilon - 1} \right)$$

$$= \frac{1}{n} \sum X_i$$

$$= \mu$$

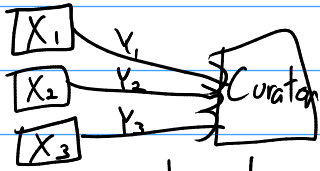
$$e^\epsilon \approx 1 + \epsilon$$

$$\text{Var}[\hat{\mu}] = \text{Var} \left[ \frac{1}{n} \sum \frac{2+\epsilon}{\epsilon} \cdot Y_i \right] = \frac{1}{n^2} \cdot n \cdot \left( \frac{2+\epsilon}{\epsilon} \right)^2 \cdot \text{Var}(Y_i)$$

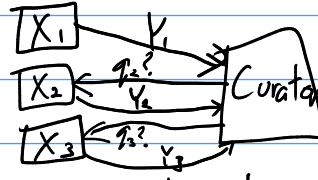
$$= O\left(\frac{1}{n \epsilon^2}\right)$$

Central DP:  $|\hat{\mu} - \mu| \leq \frac{1}{\epsilon n} \Leftrightarrow$  If want  $|\hat{\mu} - \mu| \leq \alpha$ ,  $n \geq \frac{1}{\alpha \epsilon}$ .

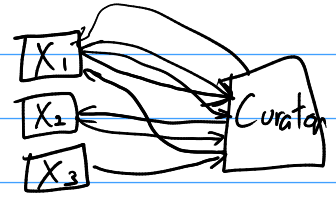
LDP:  $|\hat{\mu} - \mu| \leq \frac{1}{\epsilon \sqrt{n}}$ .  $\Leftrightarrow |\hat{\mu} - \mu| \leq \alpha$  reqs  $n \geq \frac{1}{\alpha^2 \epsilon^2}$



Non interactive



Seq. Interactive



Interactive

$\epsilon$ -LDP  $\approx$   $(\epsilon_1 \delta)$ -LDP?

Warner '65

Efimievski, Gehrke, Srikant '03

Kasiviswanathan, Lee, Nissim, Raz Khadnikova, Smith '08

KLeeNRS

Duchi, Jordan, Wainwright '13

# Local DP at Google

RAPPOR - Erligsson, Pihur, Korolova '14

$$\epsilon = \ln 3 \Rightarrow RR = \begin{cases} \text{truth w.p. } \frac{3}{4} \\ \text{lie w.p. } \frac{1}{4} \end{cases}$$

RR = yes w.p.  $\frac{1}{4}$  if user's value is no  
100 days in a row

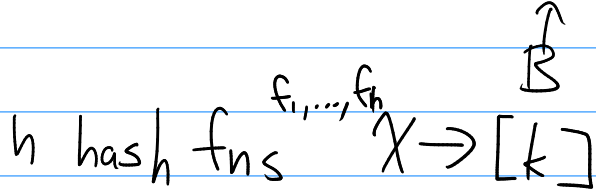
RR = 75 yes's, 25 no's. w.p.  $1.39 \times 10^{-24}$  if val = no

$$\epsilon = \ln 3 \approx 1.1 \Rightarrow 100\epsilon \approx 110, \quad e^{110}$$

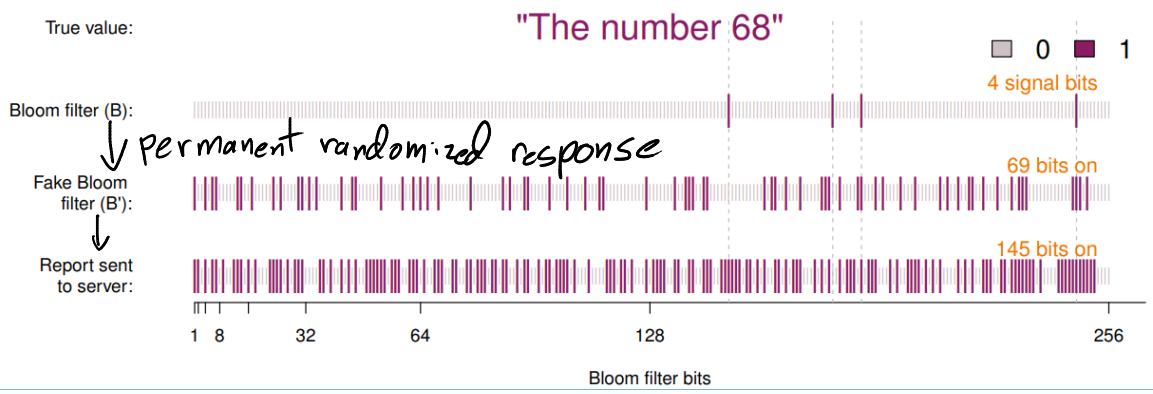
Memorization: Don't re-randomize <sup>ans to</sup> same q.

RR 0,1

RAPPOR  $v \in X \rightarrow \{0,1\}^k$       $B_j = 1$  if  $\exists i$  st  $f_i(v) = j$



~~$B = \{1, \dots, k\}$~~   
 $h=1 \rightarrow$  one-hot encoding



10000, 10001, 10002

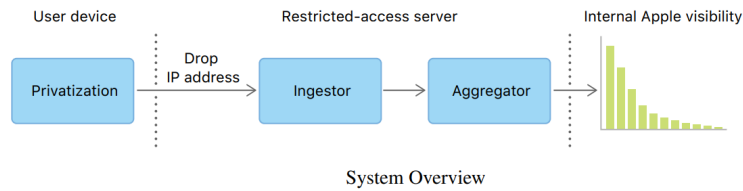
↑            ↑            ↑

$\epsilon$              $\epsilon$              $\epsilon$

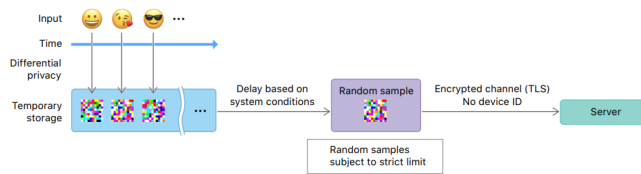
Chrome  
- malware  
- process

# Local DP at Apple '17

Tang et al. '17



System Overview



Privatization Stage

# Charikar, Chen, Farach-Colton '02

## Communication.

aaaaa aaaaa  $\rightarrow$  [00 100000 1 1 0 00]

aaaaa aaaaab  $\rightarrow$

c  $\rightarrow$   
d  $\rightarrow$

## Sequence Fragment Puzzle

1. Discover new words
2. Compute frags of words

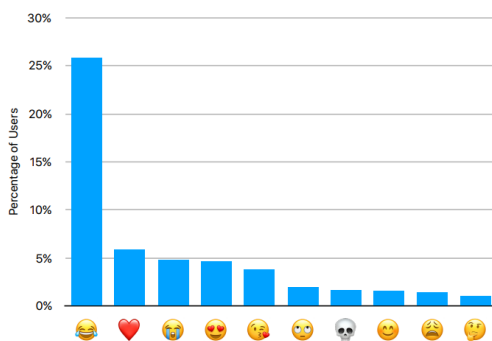
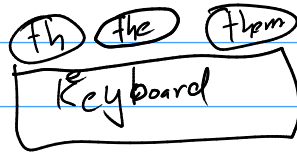
Waterloo.  $f: \text{word} \rightarrow \{0, \dots, 255\}$

$f(\text{Waterloo}) = 42.$

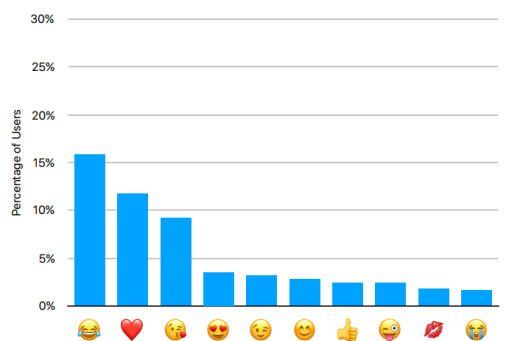
(1, wa, 42) (3, te, 42) (5, rl, 42) (7, oo, 42)

Waterloo.

iPhone Messages  
wyd wbu idc  
bruh bae  
lov , th



(a) English



(b) French

# Differential Privacy at Microsoft

Ding, Kulkarni, Yekhanin '17

## Communication

Average ~~20,13~~  $[0, m]$   $\epsilon$ -LDP

$$X_i \in [0, m] \rightarrow Y_i = X_i + \text{Lap}\left(\frac{m}{\epsilon}\right)$$

$$\hat{\mu} = \frac{1}{n} \sum (X_i + \text{Lap}\left(\frac{m}{\epsilon}\right))$$

$$\hat{\mu} - \mu = \frac{1}{n} \sum \text{Lap}\left(\frac{m}{\epsilon}\right)$$

$$\text{mag} \approx \frac{m}{\epsilon \sqrt{n}}. \quad n \geq \frac{m^2}{\alpha^2 \epsilon^2} \rightarrow \text{err} \leq \alpha.$$

Optimal

Comm:  $\Omega(\log m)$  bits

1 bit protocol.

$$X_i \in [0, m]. \quad Y_i = \begin{cases} 1 & \text{w.p. } \frac{1}{e^\epsilon + 1} + \frac{X_i}{m} \cdot \frac{e^\epsilon - 1}{e^\epsilon + 1} \\ 0 & \text{o.w.} \end{cases}$$

$$\hat{\mu} = \frac{m}{n} \sum \frac{Y_i (e^\epsilon + 1) - 1}{e^\epsilon - 1}$$

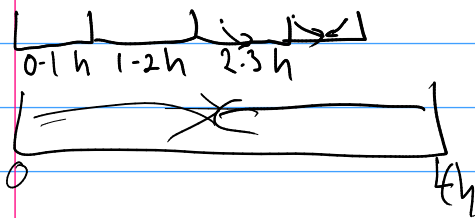
$$E[\hat{\mu}] = \frac{m}{n} \sum \frac{\left(\frac{1}{e^\epsilon + 1} + \frac{X_i}{m} \cdot \frac{e^\epsilon - 1}{e^\epsilon + 1}\right) (e^\epsilon + 1) - 1}{e^\epsilon - 1} = \frac{m}{n} \sum \frac{\left(1 + \frac{X_i (e^\epsilon - 1)}{m}\right) - 1}{e^\epsilon - 1}$$

$$= \frac{1}{n} \sum X_i = \mu$$

Chernoff-Hoeffding  $\Rightarrow \text{err} \leq O\left(\frac{m}{\epsilon \sqrt{n}}\right) \Leftrightarrow n \geq \frac{m^2}{\alpha^2 \epsilon^2} \rightarrow \text{err} \leq \alpha.$   
but only 1 bits

# Memoization

- Doesn't handle small val changes



$$X_i \in [0, m]$$

1. Memoize: Precompute  $A_i(0)$  and  $A_i(m)$
2. Select offset:  $\alpha_i \in \{0, \dots, m-1\}$  v.a.r.
3. If  $X_i + \alpha_i \leq m$ , report  $A_i(0)$   
else,  $A_i(m)$

$$A_i(0), A_i(0), A_i(0), \dots, A_i(m), A_i(m), \dots$$